

# Parallel Behavior Composition for Manufacturing

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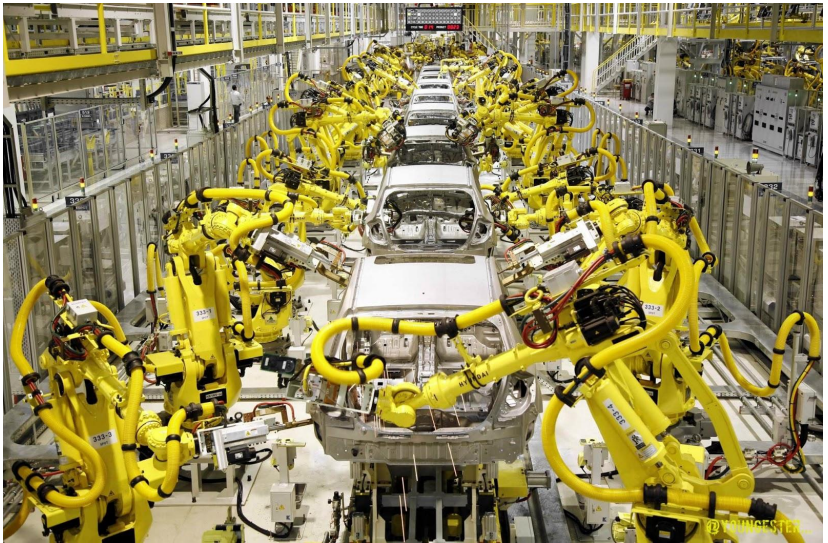
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# The problem

- manufacturers increasingly need to produce variable volumes of highly customised products, rapidly and at low cost
- one way to do this is to allow production control software greater autonomy in determining how products will be manufactured
- requires the *automated synthesis of controllers that are able to manufacture any instance of a given product type on a particular production or assembly line*
- **problem:** standard AI behavior composition, e.g., (De Giacomo et al 2013), assumes a **single sequential execution** of the target behavior module



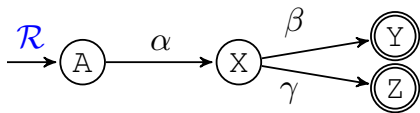
want to manufacture **multiple instances of a product in parallel**

# Key contributions

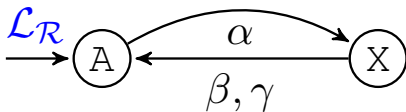
- we extend classical AI behavior composition to manufacturing settings
- we introduce a novel solution concept for manufacturing composition, **target production processes**, that are able to manufacture **multiple instances of a product simultaneously** in a given production plant
- we propose a technique for synthesizing the **largest target production process**, together with an associated **controller** for the production plant

# Production recipes & cycles

- a **production recipe** is the (finite) sequence of steps required to manufacture any instance (e.g., a blue shirt of size XL) of a particular product type (e.g., shirts)

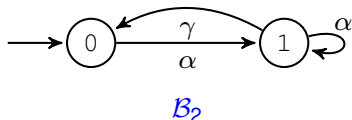
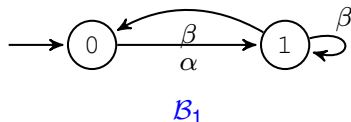


- the **production cycle**  $L_{\mathcal{R}}$  induced by a recipe  $\mathcal{R}$  represents the repeated execution of  $\mathcal{R}$



## Production plants & controllers

- production recipes are enacted by a production plant, consisting of *production resources* and their associated *capabilities*
- the capabilities of each resource are modelled as an **available behavior** — a (possibly nondeterministic) transition system



- a **production plant system**  $\mathcal{S}$  consisting of  $m$  resources is the synchronous product of  $m$  available behaviors (plus “no-op” ‘-’ actions)
- a **plant controller** is a function which delegates  $k \leq m$  actions to  $m$  available behaviors in the plant

# Target production process

- production cycle  $\mathcal{L}_{\mathcal{R}}$  tells us how to make items, and the available behaviors tell us what the production plant can do
- but the **target production process** specifying how to make multiple items of a product type simultaneously is not given
  - note that the target production process is *not* just the synchronous execution of multiple copies of  $\mathcal{L}_{\mathcal{R}}$
  - the production plant may be able to make  $n$  instances of an item in parallel, but in *only one* (or a small number of) way(s)
- we need a **new solution concept** for manufacturing composition

# TPP requirements

- R1.** a TPP should allow the manufacture of multiple product items simultaneously
- R2.** for each product item in a TPP the complete product recipe should be available
- R3.** the evolution of a TPP may depend on how (nondeterministic) behaviors in the plant happen to evolve
- R4.** a TPP may “pause” certain items (using “no-op” actions)
- R5.** a TPP should never allow the starvation of any item



# Production controllers

**$n$ -pcycle**  $\mathcal{L}^n$  the **synchronous execution** of  $n$  copies of the production cycle  $\mathcal{L}$

achieves **R1**, but is *too demanding*: represents *all possible ways* in which  $n$  items can be manufactured

**$n$ -product**  $\mathcal{T}$  fragment of a  $n$ -pcycle with **memory & uncontrollable transitions**

captures that not every interleaving of a  $n$ -pcycle may be possible in the production plant

**$n$ -TTP** an  $n$ -product that is **complete** and **fair** w.r.t. the  $n$ -pcycle

Given a recipe  $\mathcal{R}$  and production plant system  $\mathcal{S}$

A **production controller** is a pair  $\langle \mathcal{T}, P \rangle$ , such that  $\mathcal{T}$  is an  **$m$ -TTP** for  $m$ -pcycle  $\mathcal{L}_{\mathcal{R}}^m$  and  $P$  is a **plant composition** for  $\mathcal{T}$  in  $\mathcal{S}$

# Largest TPP

- we can compute a production controller  $\langle \mathcal{T}, P \rangle$  where  $\mathcal{T}$  is the **largest realizable production process**, i.e., which can “mimic” *all* realizable production processes
- define a special notion of “**simulation**” (lazy simulation) between the production plant  $\mathcal{S}$  and the  $m$ -pcycle  $\mathcal{L}_{\mathcal{R}}^m$ , s.t.:
  - $\mathcal{S}$  “mimics”  $m$  copies of the production cycle  $\mathcal{L}_{\mathcal{R}}$
  - $\mathcal{S}$  can delegate all action vectors of  $\mathcal{L}_{\mathcal{R}}^m$ , but *not necessarily in a step-by-step fashion*: product instances may be paused
- compute the set encoding memory and the set of uncontrollable transitions in the  $m$ -product
- both can be extracted from the **enacted system**: joint execution of  $m$ -pcycle  $\mathcal{L}_{\mathcal{R}}^m$  with production plant  $\mathcal{S}$

# Main result

- we say that a state  $s$  of  $\mathcal{S}$  **lazily simulates** a state  $t$  of  $\mathcal{L}^m$ , denoted  $s \rightsquigarrow t$ , iff there exists a lazy simulation relation  $\sigma_+$  of  $\mathcal{L}^m$  by  $\mathcal{S}$  with  $\langle s, t \rangle \in \sigma_+$
- we say that  $\mathcal{S}$  lazily simulates  $\mathcal{L}^m$ , denoted  $\mathcal{S} \rightsquigarrow \mathcal{L}^m$ , iff  $s_0 \rightsquigarrow t_0$
- if  $\mathcal{S} \rightsquigarrow \mathcal{L}^m$  we can build a **maximal controller generator** (CG), that returns, at each step, the set of all possible action vectors that can be delegated to behaviors

## Theorem

Let  $\mathcal{C}$  be the maximal CG for  $\mathcal{L}^m$  and  $\mathcal{S}$ . Then (i) any  $P$  is a plant composition for  $\mathcal{T}$  in  $\mathcal{S}$  iff  $\langle \mathcal{T}, P \rangle$  is generated by  $\mathcal{C}$  and (ii)  $\mathcal{T}$  is the largest realizable  $m$ -TPP for  $\mathcal{L}$  in  $\mathcal{S}$ .

## Summary & future work

- we have defined a notion of “adequacy” for solutions in the form of TPPs that respect requirements **R1-R5**
- in future work we plan to *refine* *m*-TPPs to give “efficient” manufacturing processes, e.g., with respect to average throughput, machine utilization, load balancing, etc.
- such optimizations can be done after the TPP has been built, or possibly during synthesis by discriminating between controllers

For more details, please see poster ??